

## *A Printer Turing Machine*

A **printer-TM** is TM with an added printer-tape. The printer-TM writes strings on the printer-tape (separated by  $\Delta$ ); once written, a string is not altered.

**Theorem.** *A language is r.e. if and only if some printer-TM outputs precisely those strings.*

The proof is in two constructions...

## *From Printer-TM to Standard TM*

Armed with printer-TM  $M$  for language  $L$ , we build standard TM  $N$ .

On input  $x$ , TM  $N$  runs  $M$  and monitors  $M$ ; if  $N$  ever finds  $x$  on the printer-tape, then  $N$  accepts. So  $N$  accepts strings in  $L$ , and does not halt otherwise.

## *From Standard TM to Printer-TM*

Armed with standard TM  $N$  for  $L$ , we build printer-TM  $M$ .

The idea is to run  $N$  on every possible string in parallel—an infinite number of tasks!—and print out those it accepts.

## *Infinite Parallelism*

The printer-TM works  $M$  in rounds.

In round  $i$ ,  $M$  starting from scratch, generates the first  $i$  strings lexicographically (in dictionary order), runs  $N$  on each for  $i$  steps, and outputs any string that is accepted.

Eventually, every string in  $L(N)$  will be generated and  $N$  run for long enough, and will appear in the output.