

Regular Grammar

This is grammar where every rule has right-hand terminal OR terminal variable

Example:

$$A \rightarrow 1A \mid 0B$$

$$B \rightarrow 1A \mid 0C \mid 0$$

$$C \rightarrow 0C \mid 1C \mid \emptyset \mid 1$$

There is algorithm to convert DFA to regular grammar.

state $\Rightarrow$ variable

transition $\Rightarrow$ production
with terminal
& variable

transition
to accept $\Rightarrow$ production with
terminal

