

# The Internet Protocol

- IP Addresses
- Address Resolution Protocol:
- IP datagram format and forwarding:
- IP fragmentation and reassembly

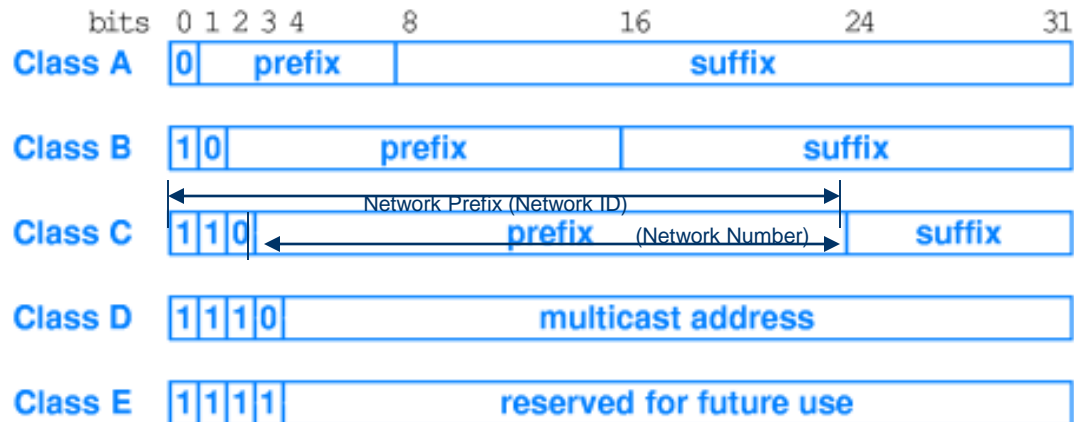
# IP Addresses

- IP Addresses are 32 bit.
- Written in dotted decimal format: X.X.X.X
  - Example: A Clemson address : 130.127.48.4
- An address encodes the identification of the network as well as the host (network id, host id)
- What is a host?
- How does an address relate to a host ?
- Three types of addresses?

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  - Example: A Clemson address : 130.127.48.4
- An address encodes the identification of the network as well as the host (network id, host id)
- What is a host? A computer that can communicate with another computer over a TCP/IP network.
- How does an address relate to a host ?
  - Not one –to-one !!
  - Multihomed hosts have more than one interface AND more than one IP address
  - So slightly better answer is one IP address per interface although....
    - NAT totally changes this - Network Address Port Translation allows any number of Hosts to share the same IP address
- Three types of addresses?
  - Unicast
  - Broadcast
  - Multicast

# Original Classful Addressing Scheme



Original address scheme was classful:

- Class A for large networks (>64K hosts)
- Class B for medium networks (>256 hosts)
- Class C for small networks (<256 hosts)
- Class D for multicast
- Class E reserved

# Original Classful Addressing Scheme

- What's the total address space?
- Example: Clemson address: 130.127.49.225
  - Class ?
  - Network ID?
  - Host ID?
- Example: 192.168.1.100

# Broadcast Addresses: Two Types

- A network directed broadcast
  - contains a valid network id with an all 1's host id
  - class C example: 192.168.1.255
  - From a host on a different network (192.168.2.2), conceptually, if we issue 'ping -c 192.168.1.255', ping should receive an echo response from all hosts on the 192.168.2 network.
    - There's not a router in the world that will allow this!!!!
- A local or limited broadcast does not require knowledge of the network address.
  - Referred to as the all 1's broadcast:  
255.255.255.255
  - A limited broadcast useful for certain startup protocols:

# Special Addresses

- Conventions.... The 'this' rule and the 'all' rule:
  - A netid or a host id of '0' implies 'this'
  - A netid or a host id of '1' implies 'all'
- Class C example: 192.168.1.0

*The 'this' rule: 'This host' on the network*

- Class C example: 192.168.1.255

*The '1's' rule: 'All' hosts on the network*

# Addresses with Special Meaning

- All 0's and all 1's: has slightly different meanings for the entire address, just the network id or just the host id
- Loop Back : 127.x.x.x (e.g., 127.0.0.1)
  - What address class?
  - Conceptually, all  $2^{24}$  addresses in this class “A” network are equivalent.
    - In reality, the Host OS probably will only accept a few. Try
      - Ping 127.0.0.1
      - Ping 127.22.22.22
      - Ping localhost //A defacto standard host name that gets resolved to 127.0.0.1
- Private address space:
  - RFC 1918 defines certain address ranges for private use.
    - 10.0.0.0 - 10.255.255.255 (Class A space)
    - 172.16.0.0 - 172.31.255.255(Class B space)
    - 192.168.0.0 - 192.168.255.255 (Class C space)

# Special Addresses

So, how many valid host ids are available with the following address:

192.1.1.0/24

What happens if you ping 0.0.0.0 ?

# Classful Address Ranges

Class	Lowest Address	Highest Address
A	1.0.0.0	126.0.0.0
B	128.1.0.0	191.255.0.0
C	192.0.1.0	223.255.255.0
D	224.0.0.0	239.255.255.255
E	240.0.0.0	247.255.255.255

Dotted Decimal Class Address Ranges

# Classful Addresses

- Addresses do not specify computers, but rather connections to particular hosts.
- Multihomed: A host that has  $>1$  physical connection.

# Classful Addresses

Basic mechanism: A two level class hierarchy.

- requires a unique network prefix for each physical interface.
- Two additional schemes designed to conserve net addresses: subnet and classless addressing (CIDR).

Other issues:

- Mobility
- Flexibility
- Naming
- And of course the IP Address shortage problem...

# Subnetworks

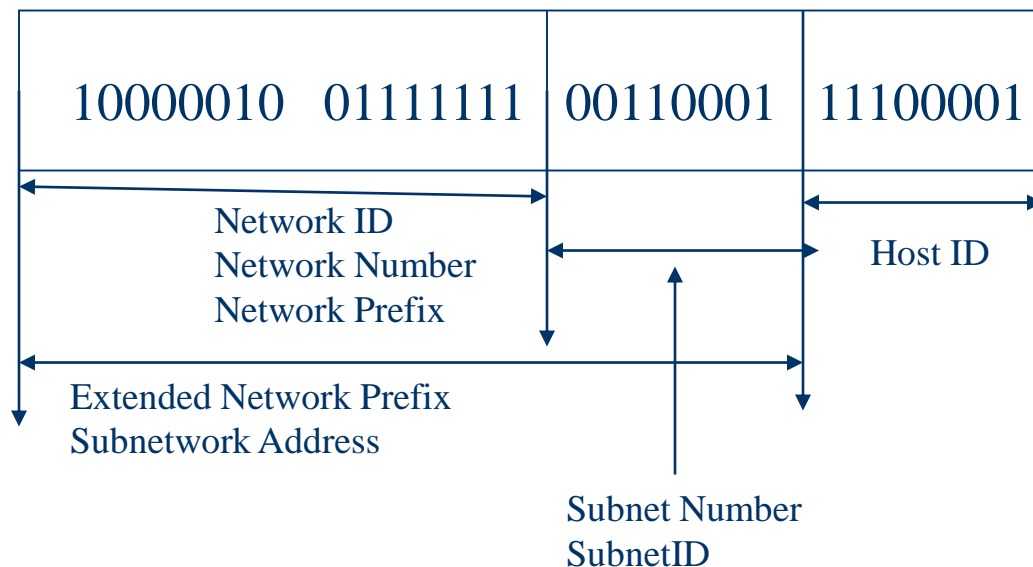
**Subnet addressing:** Allows an organization to create multiple logical networks internally while preserving the presence of a single autonomous network to the outside.

An autonomous network:

*A collection of routers and networks that fall under one administrative entity.*

# Subnetworks

Example : 130.127.49.225 mask 255.255.255.0



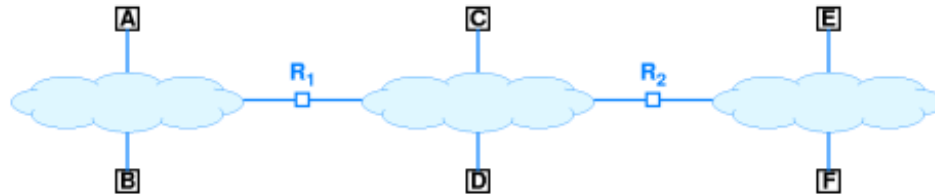
Classfull network address: 130.127.0.0

Classfull network prefix notation: 130.127/16

Subnetwork address: 130.127.49.0

Network Prefix Notation: 130.127.49/24

# Address Resolution Protocol



The Address Resolution Problem:  
 In order for two hosts on the same network to communicate, they need to know each others MAC address.



Dest Address (6)	Source Address (6)	Length (2)	DSAP (1)	SSAP (1)	cntl (1)	code (3)	type (2)	Data ( 38 - 1492)	CRC (4)
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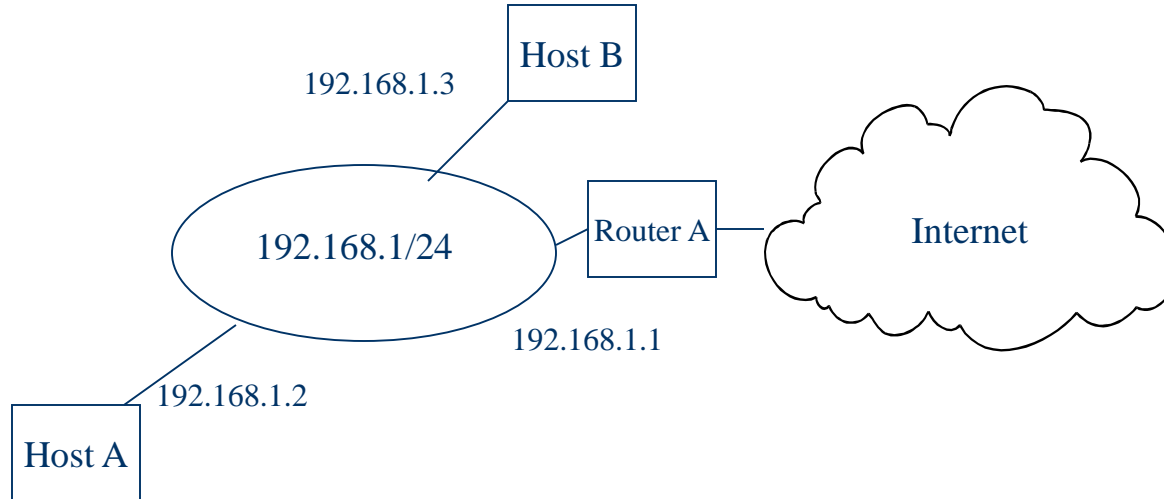
# Address Resolution Protocol

<u>IP Address</u>	<u>Hardware Address</u>
197.15.3.2	0A:07:4B:12:82:36
197.15.3.3	0A:9C:28:71:32:8D
197.15.3.4	0A:11:C3:68:01:99
197.15.3.5	0A:74:59:32:CC:1F
197.15.3.6	0A:04:BC:00:03:28
197.15.3.7	0A:77:81:0E:52:FA

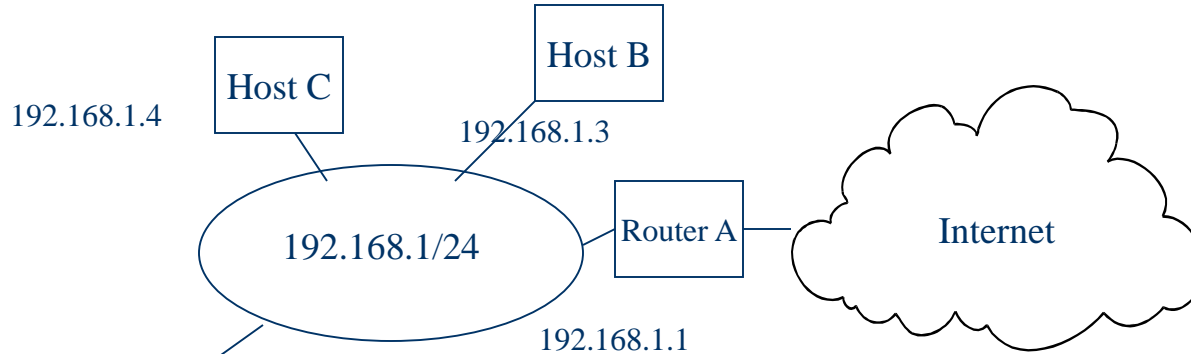
One approach is through direct binding:  $\text{Phys} = f(\text{IP})$

- Objective is to encode the IP address in Physical
- Dynamic binding protocol is required for Ethernet networks where:
  - Physical address larger than IP address (48 vs 32 bits)
  - Physical address can change
  - Potentially many hosts on the network
- Address resolution protocol (ARP) is a table lookup approach
  - Defined by RFC 826

# ARP: Host A Sends to Host B



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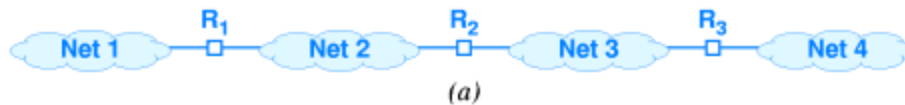
Host A's arp cache:

IP ADDR	MAC ADDR
192.168.1.4	xxxxxx

- Host A looks in its ARP table to see if the binding exists
- If not, Host A issues a limited broadcast (all 1's) sending an ARP 'whois' message.
- All Hosts on the network receive the broadcast. Remember that on a LAN, the frames that are received must have a matching destination MAC address or the dest.Address must be the Ethernet broadcast address.
- Host C ignores the broadcast, Host B receives the message and passes it to its ARP program
- Host B replies with a response that says map my IP address to this MAC addr

**Try an arp -a on a windows machine!!**

# The Internet Protocol (IP)



Destination	Next Hop
net 1	R <sub>1</sub>
net 2	deliver direct
net 3	deliver direct
net 4	R <sub>3</sub>

(b)

(a) An example internet with three routers connecting four physical networks, and (b) the conceptual routing table found in router R<sub>2</sub>. Each entry in the table lists a destination network and the next hop along a route to that network.

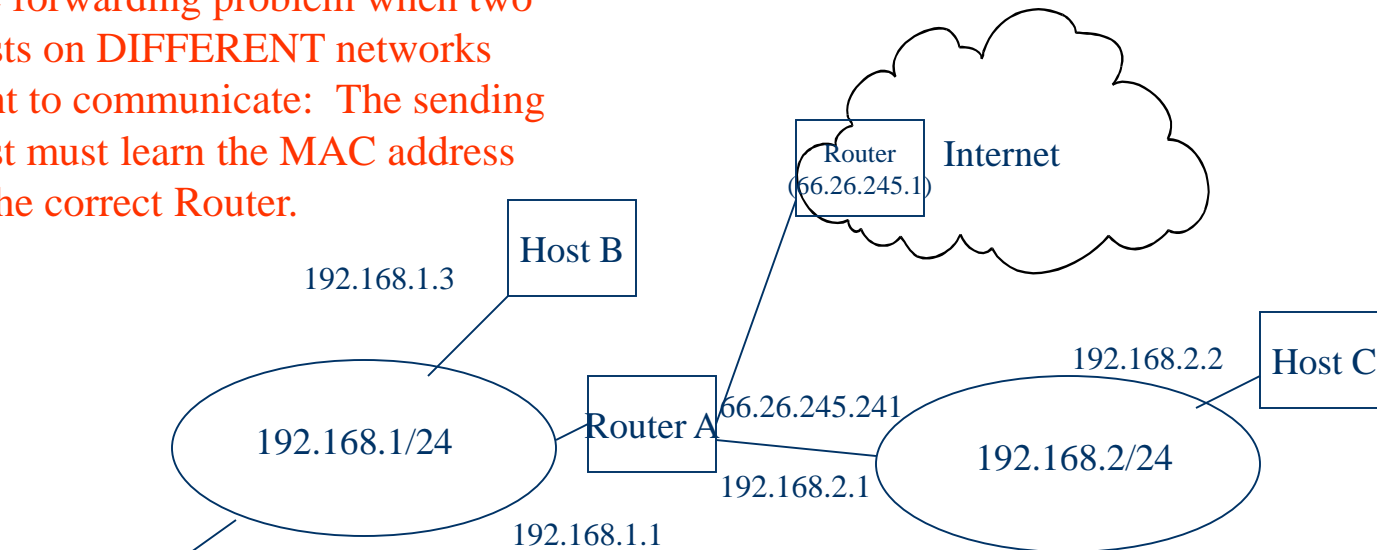
- The IP standard (RFC 791) defines:
  - The basic unit of transfer (the datagram)
  - Routing/data forwarding
  - Rules for how hosts/routers should deal with packets and errors

# Routing/Forwarding

- **Routing** means process of selecting a path over which to send packets.
- **Router** is the computer which makes such a selection.
  - The Internet is composed of multiple physical networks inter-connected by computers called *routers*.
- **Direct Delivery** - transmission of a datagram from one machine across a physical network directly to another.
- **Indirect Delivery** - when the destination is not on a directly attached network; sender passes the datagram to a router for delivery.

# Indirect Delivery: Host A sends to Host C

The forwarding problem when two Hosts on DIFFERENT networks want to communicate: The sending Host must learn the MAC address of the correct Router.



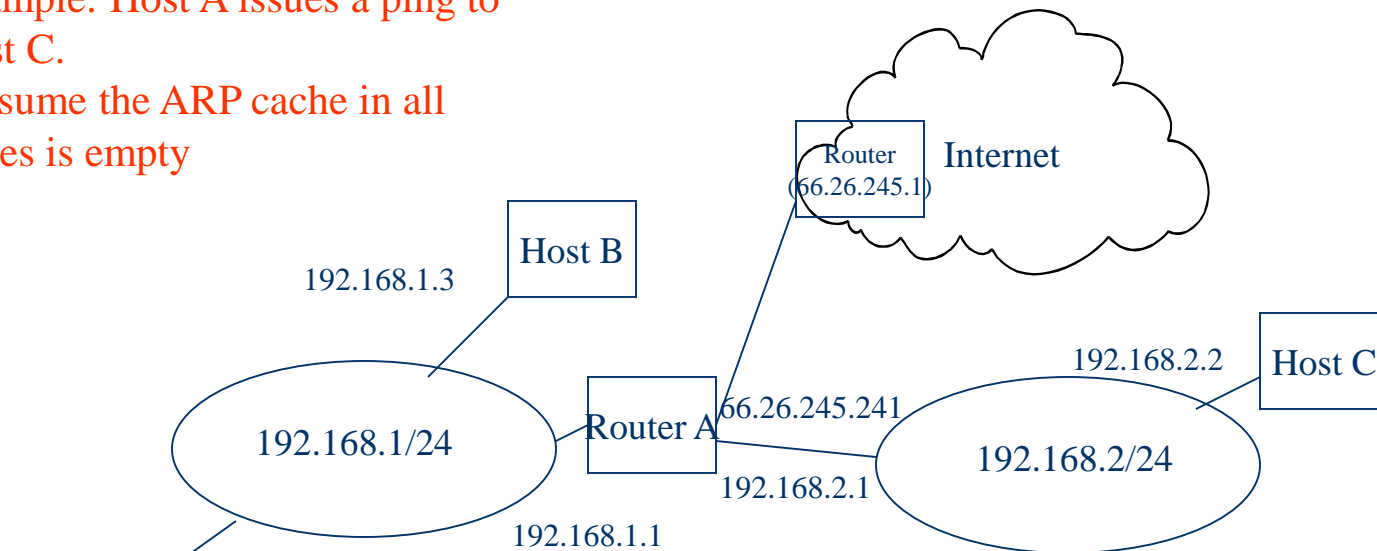
Destination Network	Router
192.168.1.0/24	192.168.1.1 (direct)
192.168.2.0/24	192.168.2.1 (direct)
66.26.245.1	66.26.245.241 (direct)
default	66.26.245.1 (indirect)

**Try a netstat -r on a windows machine!!**

# Indirect Delivery: Host A sends to Host C

Example: Host A issues a ping to Host C.

- Assume the ARP cache in all nodes is empty



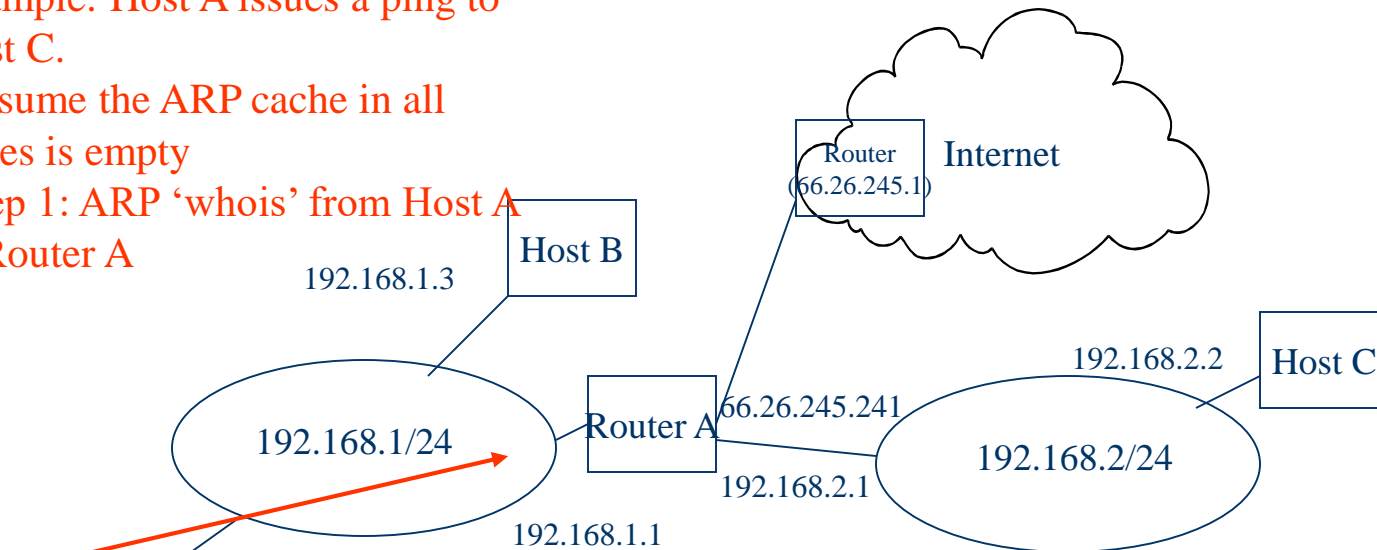
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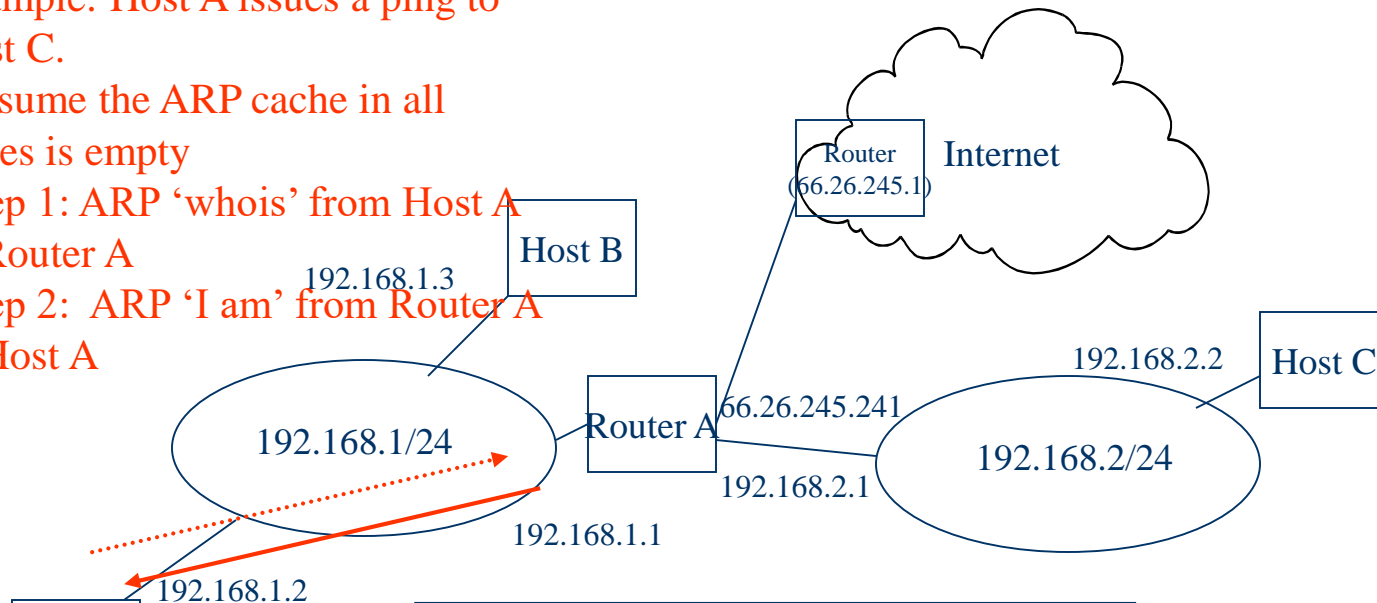
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# Indirect Delivery: Host A sends to Host C

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- Step 2: ARP 'I am' from Router A to Host A

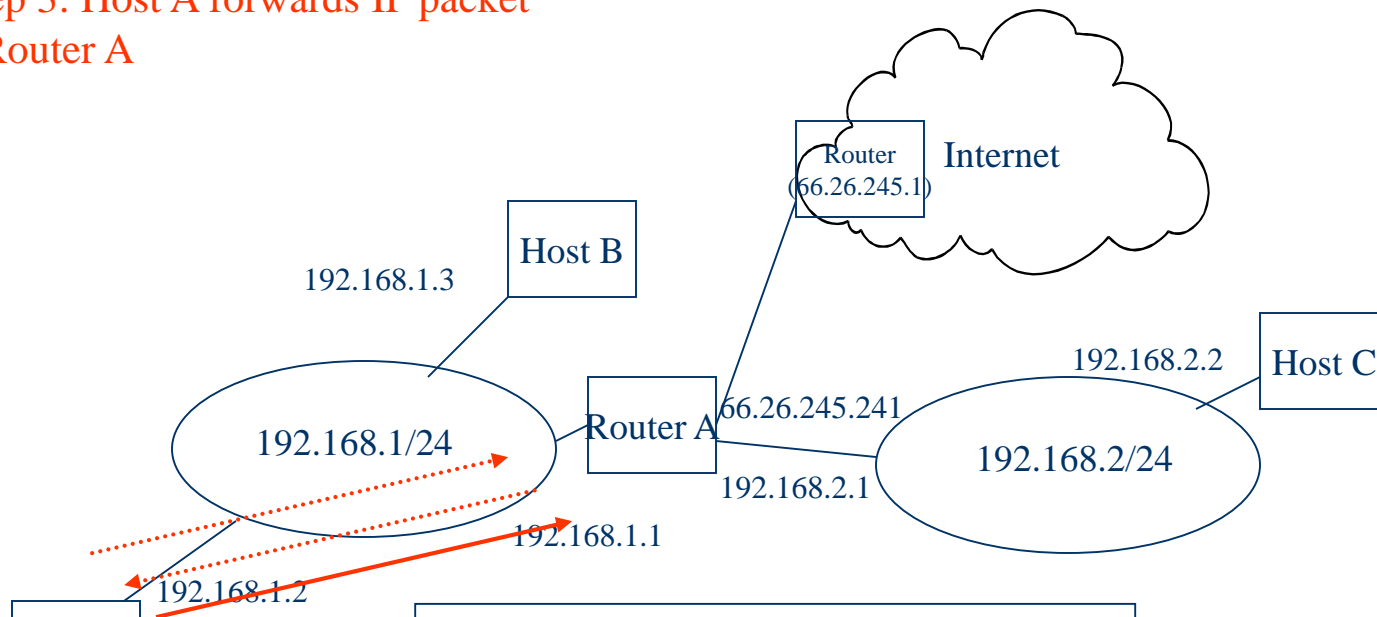


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# Indirect Delivery: Host A sends to Host C

- Step 3: Host A forwards IP packet to Router A

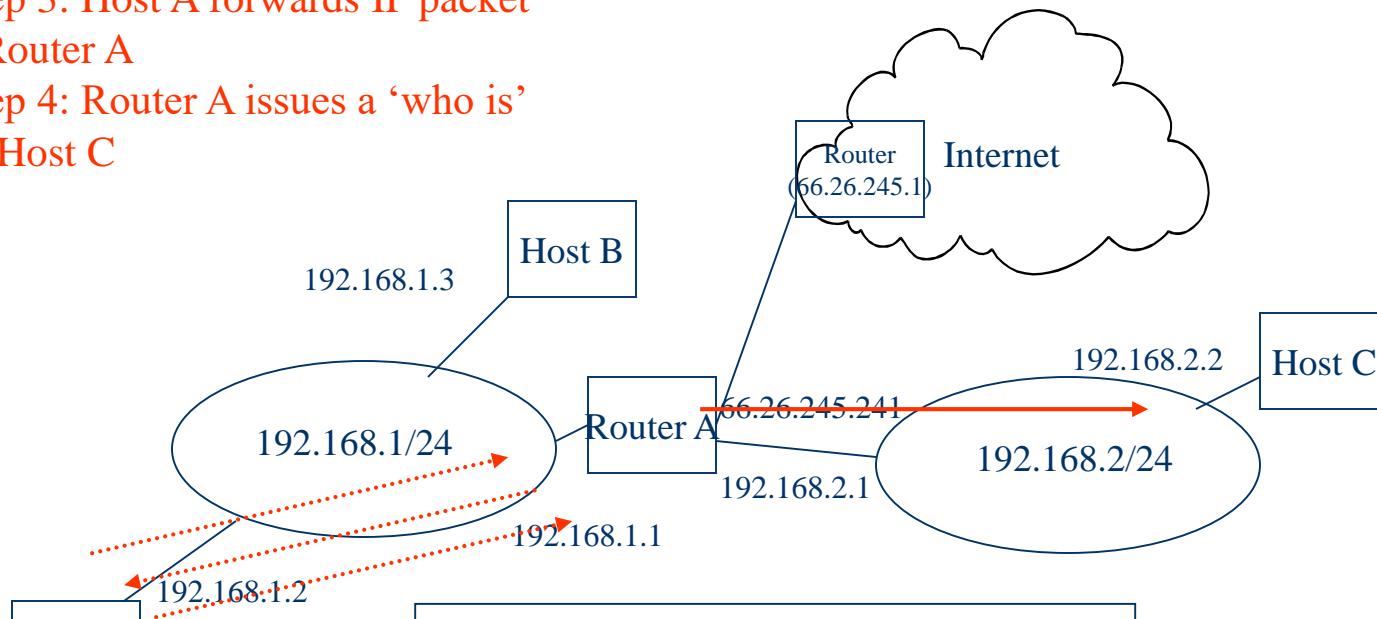


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# Indirect Delivery: Host A sends to Host C

- Step 3: Host A forwards IP packet to Router A
- Step 4: Router A issues a 'who is' for Host C

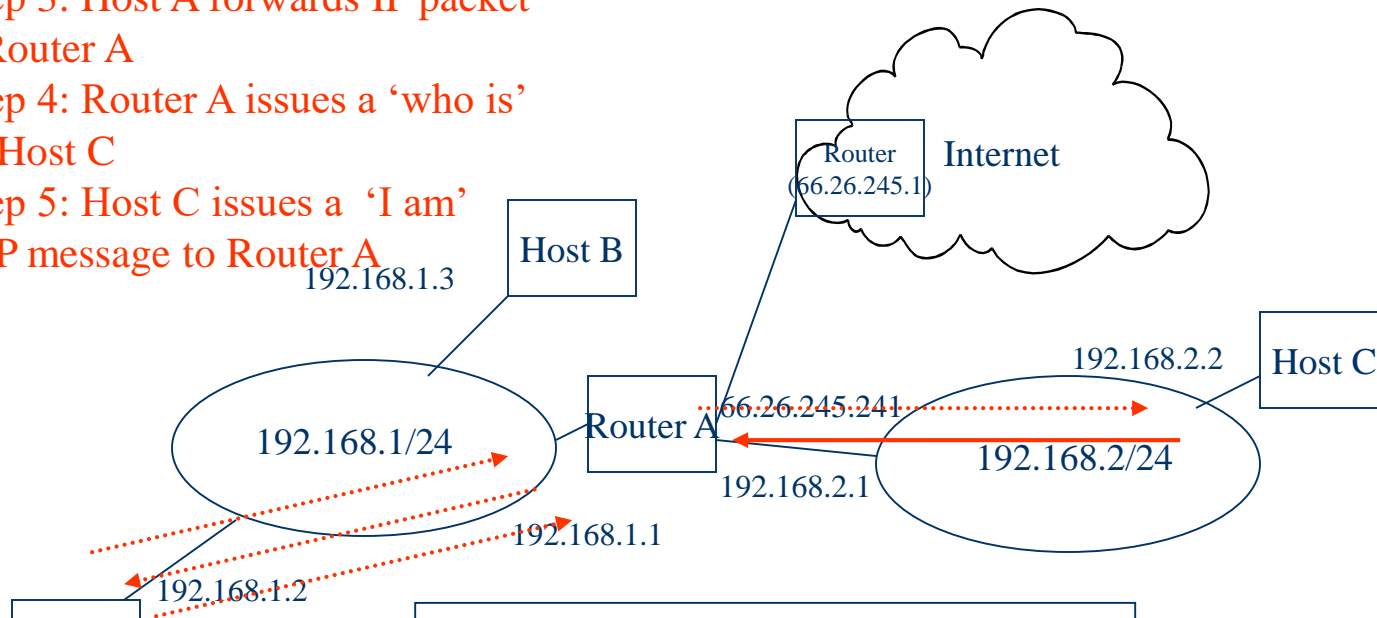


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- Step 5: Host C issues a 'I am' ARP message to Router A

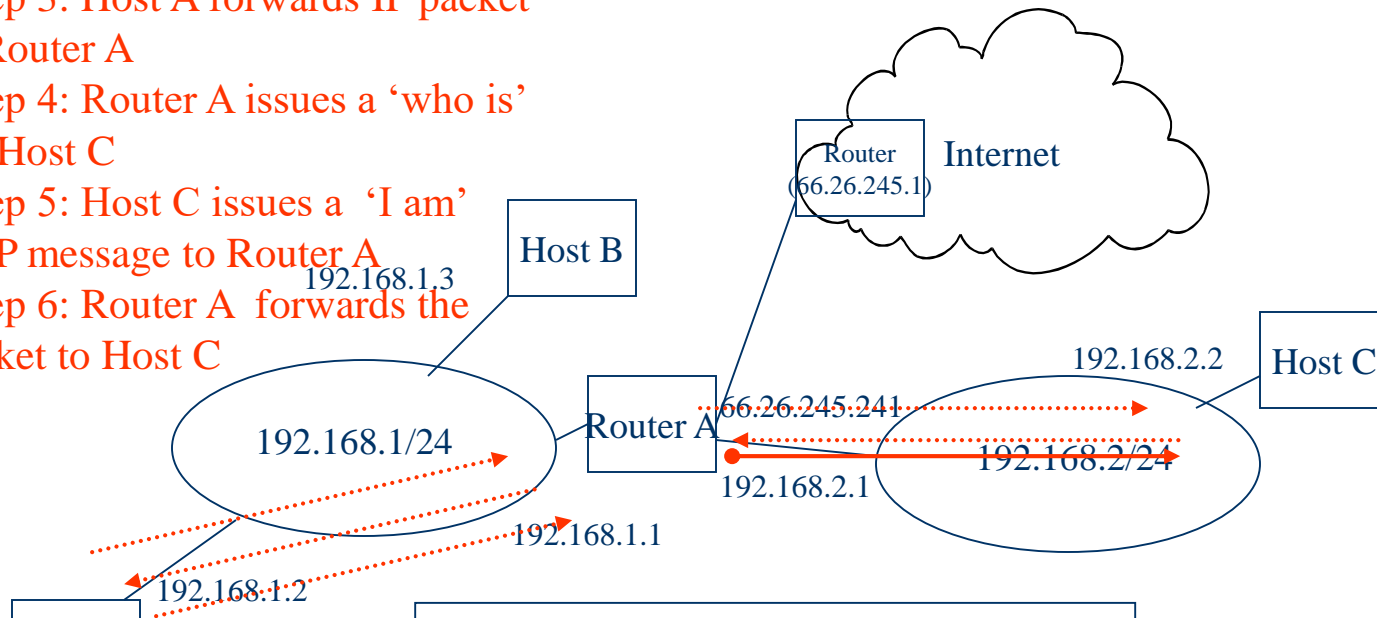


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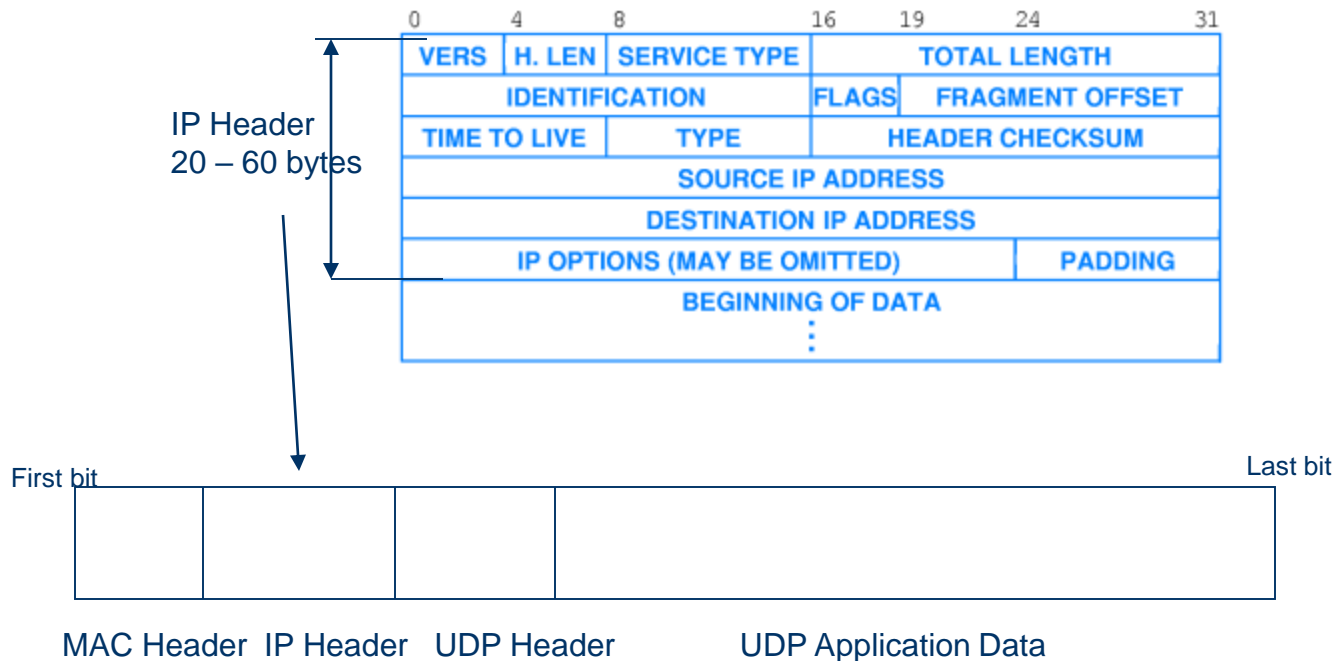
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- Step 6: Router A forwards the packet to Host C



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# IP Datagram Format: UDP example



This is what gets sent 'on the wire': a frame which contains an IP Packet

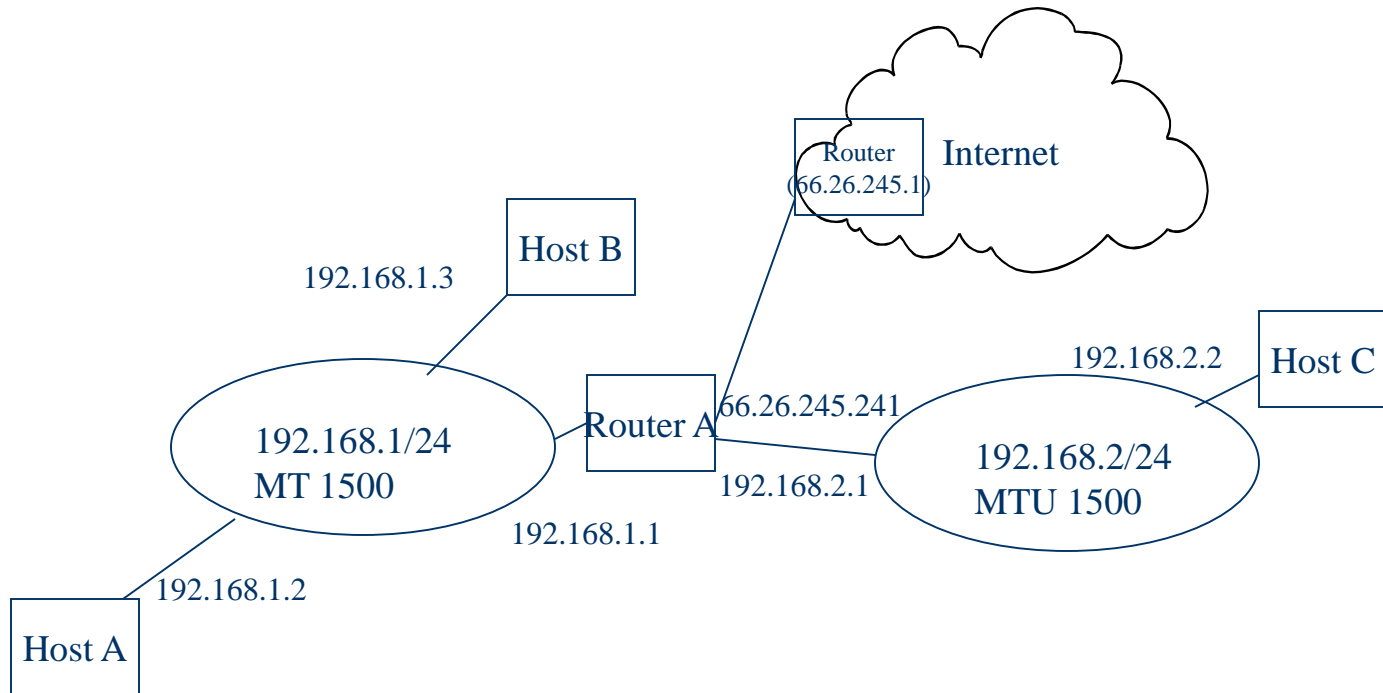
# IP Datagram Format: UDP example

Here is a 1500 byte UDP packet in hex:

```
4500 05dc  version 4, length 20bytes, TOS:0, total length:1500bytes
1234 4000  id:1234 Flags(3 bits):010 in binary (DF set), offset : 0
ff11 aaaa  TTL:255, protocol: 0x11(UDP), checksum: 0xaaaa
c0a8 0164  src:192.168.1.100
ca01 0101  dst: 202.1.1.1

5555 6666  src port: 0x5555, dst port 0x6666
05c8 bbbb  UDPLength:1480, UDP checksum:0xbbbb
...
1472 bytes of application data
...
```

# IP Fragmentation: Ping example



Ping -l 3000 192.168.2.2

# IP Fragmentation control



- Identification field: uniquely ids a datagram
- Flags: DF (don't frag) and MF (more frag)
- Fragment offset: specifies the offset in the original datagram of the data being carried in the fragment specified in units of 8 octets.
- The data portion of a fragment must be a multiple of 8 (except for the last fragment).

# IP Fragmentation: Ping example

