

## Computer Science 102 Lab 13

### Efficient methods of searching -

Suppose an online telephone book is comprised of structures of the following form:

```
struct entry_type
{
    char name[60];    // e.g. Westall, James M.
    char num[10];    // e.g. 8645551212
};
```

and it consists of  $1048576 = 1024 \times 1024 = 2^{20}$  entries as shown

```
struct entry_type phone_book[1024 * 1024];
```

Now suppose your mission is to write the function

```
struct entry_type *lookup(char *target_name);
```

The naive approach to doing this is:

```
struct entry_type *lookup(char *target_name);
{
    int i;
    for (i = 0; i < (1024 * 1024); i++)
    {
        if (strcmp(target_name, phone_book[i]) == 0)
            return(&phone_book[i]);
    }
}
```

For this approach, the length of the search (number of array elements that must be compared against target\_name) is:

- best case: 1 (You were asked to lookup the first element)
- worst case:  $1024 * 1024$  (You were asked to lookup the last element)
- average case:  $1024 * 1024 / 2$

If the phonebook is in random order, this is the best possible result.

However if the phonebook is sorted by name (as a real phonebook is) a more clever algorithm can produce the following results.

- best case: 1 you found the element on the first look
- worst case: 20 (instead of  $1024 * 1024$ )
- average case: left as an exercise (but obviously between 1 and 20)

This better algorithm is known as binary search and belongs to the class of algorithms known as *divide and conquer*. It begins by comparing the "middle element", `phone_book[1024 * 1024 / 2].name` with `target_name`.

If `target_name > phone_book[1024 * 1024 / 2].name` then the whole first half of the phone book is known NOT to contain the target name. The next step is to compare `target_name` with `phone_book[3 * 1024 * 1024 / 4].name`. At each step, the remaining range is divided by 2 and the next probe point is set to the middle of the remaining range.

Because the range of elements to be searched is divided by 2 at each step, the number of elements that must be tested is bounded by  $\log_2(\text{the table size})$ . *Therefore an element may be looked up in a sorted table of 4 BILLION elements in only 32 comparisons.*

## The bisection algorithm

Given a function  $f(x)$ , a common problem in numerical computing is to find a *value  $x$  for which  $f(x) = 0$* . If

- $f(x) = 2 * x + 7$  or
- $f(x) = x^2 - 2x + 1$ ,

you can readily solve for  $x$  using algebra.

But if the function is something like

- $f(x) = x * e^x - 10 * x$ ,

there will be no algebraic solution and if the function is something like

- $f(x) = x^2 + 1$

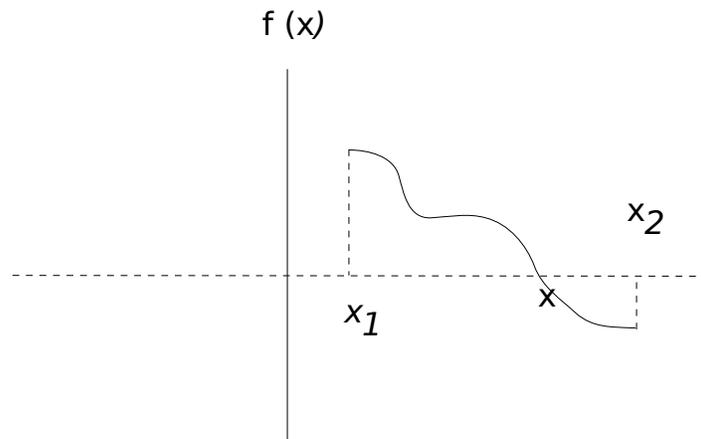
there may be NO real number solution at all!

## The *bisection* algorithm

If the function  $f(x)$  is continuous and we can find points  $x_1$  and  $x_2$  so that  $f(x_1) * f(x_2) < 0$  (the function values have different signs at  $x_1$  and  $x_2$ ), then we are guaranteed by the intermediate value theorem that there exists an  $x$  with

- $x_1 < x < x_2$  for which  $f(x) = 0$

The bisection algorithm is a divide and conquer algorithm that can be used to find the  $x$  value for which  $f(x) \sim 0$ .



```
while ((x2 - x1) > epsilon)
{
    compute xm = the midpoint between x2 and x1
    if (f(x1) * f(xm)) < 0)
        replace x2 with xm
    else
        replace x1 with xm
}
```

In this lab you will you will implement the bisect function that can find numerically a value for which an arbitrary continuous function takes on a value of approximately 0.

```
/**/  
/* This function attempts to find a zero of a function */  
/* using the bisection method */  
  
int bisect(  
double (*f)(double), // external function whose zero we seek  
double *minval, // pointer to minimum of search range  
double *maxval, // pointer to maximum of search range  
double epsilon); // tolerance
```

Your function should begin by ensuring that  $*maxval > *minval$  and that  $f(*minval)$  and  $f(*maxval)$  have different signs. If not *bisect* should return 0.

Otherwise your function should perform bisection until

$$(*maxval - *minval) < \epsilon$$

In this case the *value returned should be the number of times the interval is bisected*. Your bisection function should directly modify the values  $*minval$  and  $*maxval$  so that on return the caller can use  $*minval$  as the root of the equation  $f(x) = 0$ .

```
int main()  
{  
    int code;  
    double min;  
    double max;  
  
    min = 0.5;  
    max = 2.0;  
  
    code = bisect(f1, &min, &max, 1.0e-10);  
    printf("code = %5d min = %16.6le f(min) = %16.6le \n",  
           code, min, f1(min));
```

In this lab you will submit a single file, `bisect.c` that includes *ONLY* the new `bisect()` function constructed as part of this lab. It must not contain the `main()` function nor the `f1()` function!

```
sendlab.102.labsection# lab# bisect.c
```